

ONOS

Architecture, Abstractions & Performance



What is ONOS?



Open Network Operating System (ONOS) is an open source SDN network operating system. Our mission is to enable Service Providers to build real SDN/NFV Solutions.

ONOS Community



Brief Retrospective



- Started with a minimal platform with only a few apps
 - built with sound structure and solid code & minimalistic REST API
 - 4 apps and 1 SB plugin
- Added new core functionality and apps with each release
 - deliberately balancing investments in platform vs. use-cases and apps
 - show innovation, but also take pragmatic steps to be deployment-ready
 - maintain coherence of architecture and quality of code
- Now a platform with many features and apps
 - new capabilities, distributed primitives and even greater extensibility
 - now 70+ apps, including SB plugins, drivers, and samples

Quarterly Releases



- **Avocet** (1.0.0) released 2014-12
 - Initial release of clean and modular code-base, protocol independence
- **Blackbird** (1.1.0) released 2015-03
 - Improved performance, scale-out, increased robustness
- **Cardinal** (1.2.0) released 2015-06
 - New use-cases, additional core features, additional SB protocols
- **Drake** (1.3.0) released 2015-09
 - Platform enhancements, security, UI enhancements
- **Emu** (1.4.0) - released 2015-12
 - CORD features, prototype of dynamic cluster scaling
- **Falcon** (1.5.0) - released 2016-03
 - Dynamic cluster scaling, model extensibility, intents on flow objectives

Quarterly Releases



- **Falcon** (1.5.0) - released 2016-03
 - dynamic cluster scaling, model extensibility, intents on flow objectives
- **Goldeneye** (1.6.0) - planned for 2016-06
 - spring cleaning, intent framework, YANG tools, GUI scaling, P4 PoC
- **H...** (1.7.0) - planned for 2016-09
 - separate platform & core, network hypervisor, YANG at NB, P4 support
- ...

Platform Hardening



- Significantly improved performance
 - published white-paper and established relevant performance metrics
- Further increased quality and fault-tolerance
 - fixed defects and added a repertoire of robust distributed structures
 - fixed defects in 3rd party code and contributed changes upstream
- Improved security
 - northbound (REST, CLI & GUI), southbound and east-west secured
- Improved usability and supportability
 - deployment, component configurability, centralized app management
 - network configuration, GUI enhancements & extensibility
 - dynamic cluster scaling and model extensibility

Process Enhancements



- Established deprecation policy for API compatibility
 - give fair warning to app developers before APIs change or vanish
 - balances stability vs. ability to innovate or respond to feedback
- Incubating functionality over multiple releases
 - development of some features takes more time than a single release
 - introduce preliminary functionality in one release
 - harden & refine in the next release
- Broadening the set of code submitters
 - granting ability to +2/submit to developers based on code/review merit
 - serves both to empower the community and to off-load the core team

ONOS & Approach to SDN



- Move with urgency, but deliberately
- Mind the fundamentals & beware of yak-shaving
- Keep balance between innovation, utility and stability
- Allow *legacy* devices to participate in SDN, but not to deform or diminish the SDN vision

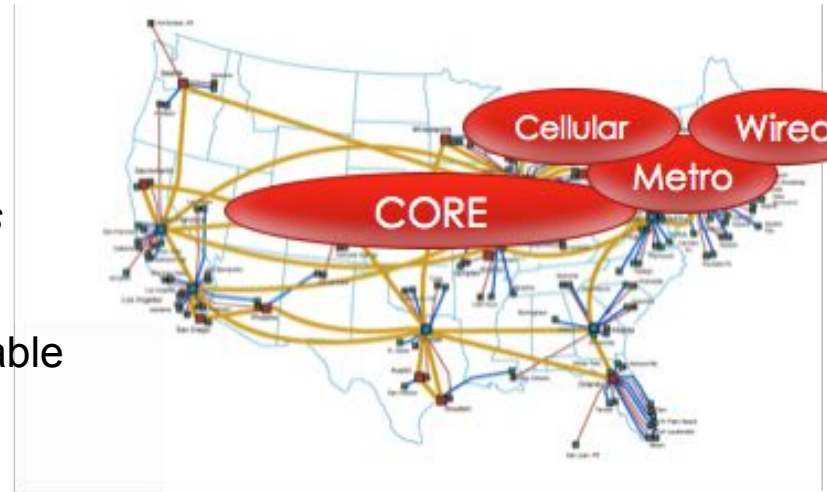


Why ONOS?

Service Provider Networks



- WAN core backbone
 - Multi-Protocol Label Switching (MPLS) with Traffic Engineering (TE)
 - *200-500 routers, 5-10K ports*
- Metro Networks
 - Metro cores for access networks
 - *10-50K routers, 2-3M ports*
- Cellular Access Networks
 - LTE for a metro area
 - *20-100K devices, 100K-100M ports*
- Wired access / aggregation
 - Access network for homes; DSL/Cable
 - *10-50K devices, 100K-1M ports*



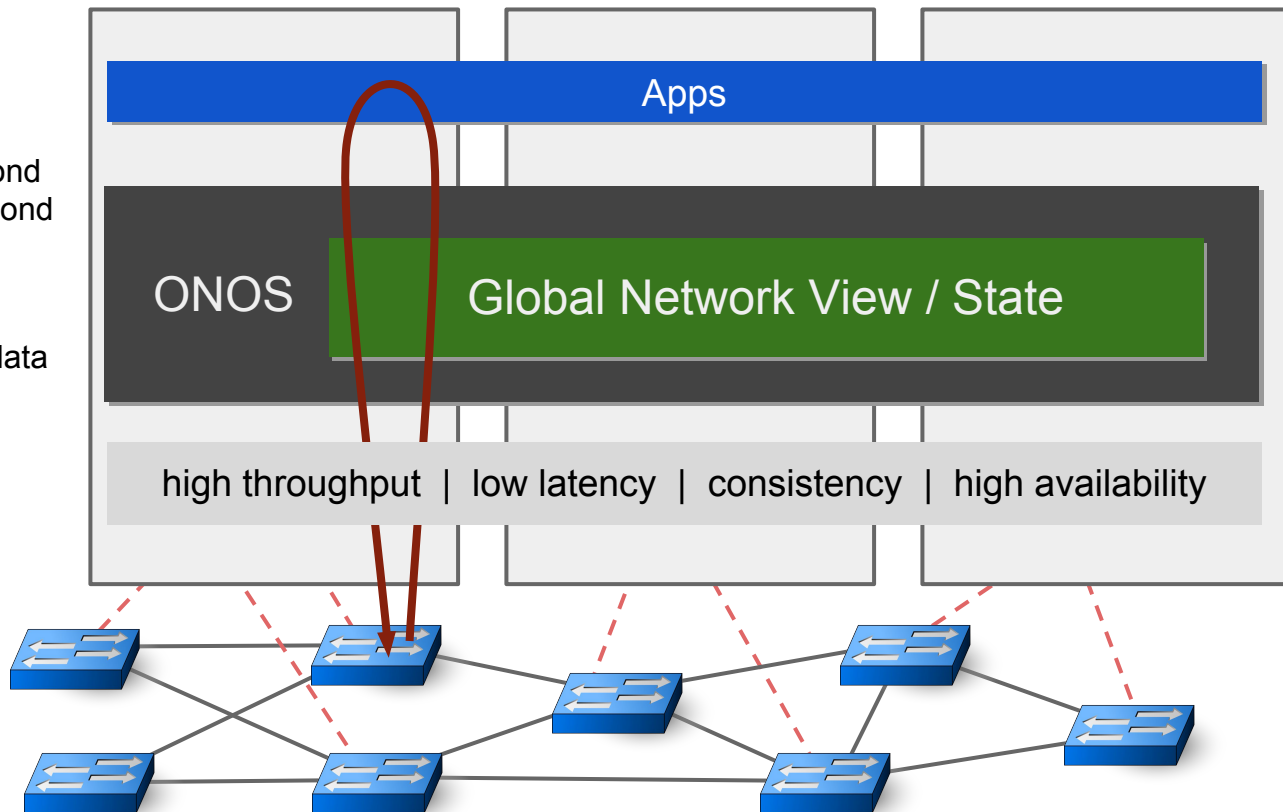
Key Performance Requirements



High Throughput:
~500K-1M paths setups / second
~3-6M network state ops / second

High Volume:
~500GB-1TB of network state data

Difficult challenge!

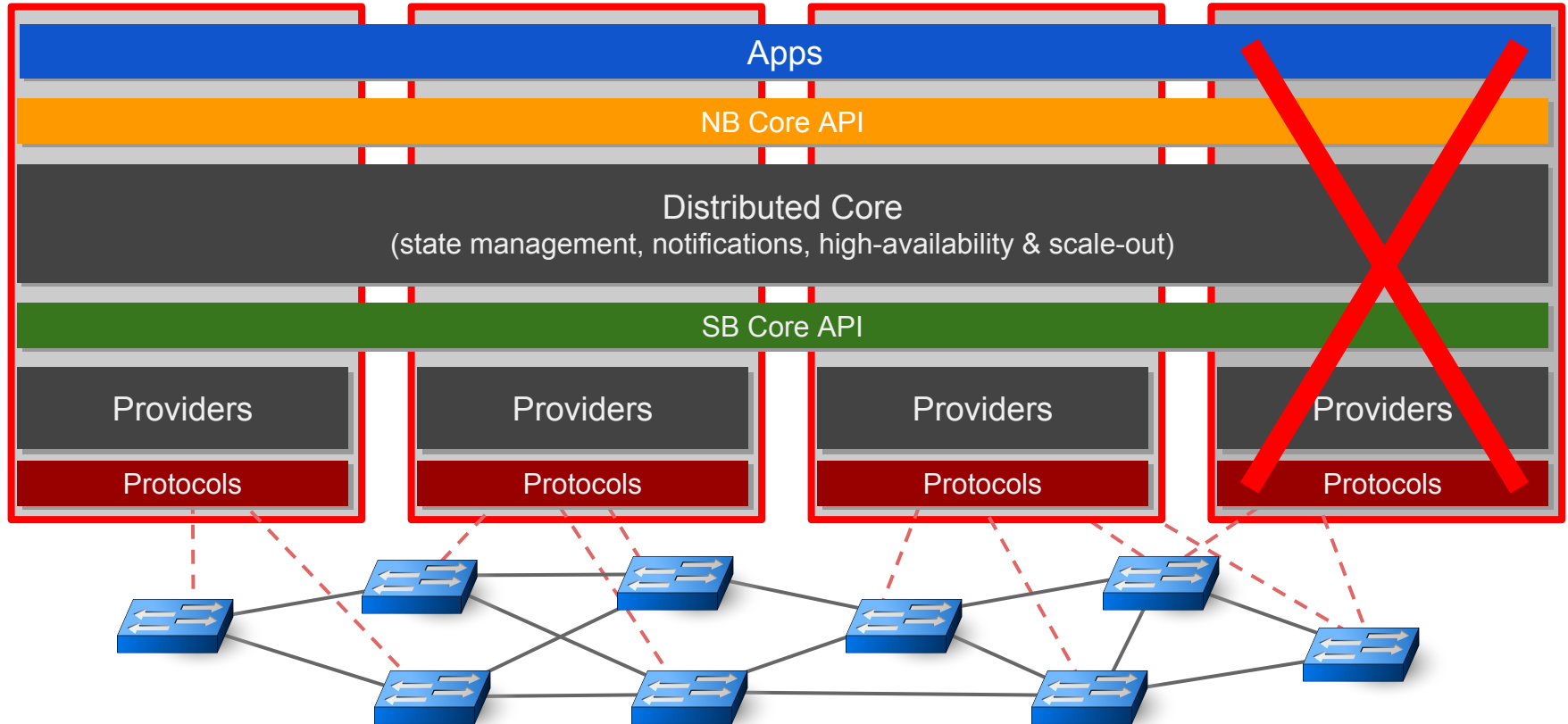


Architectural Tenets

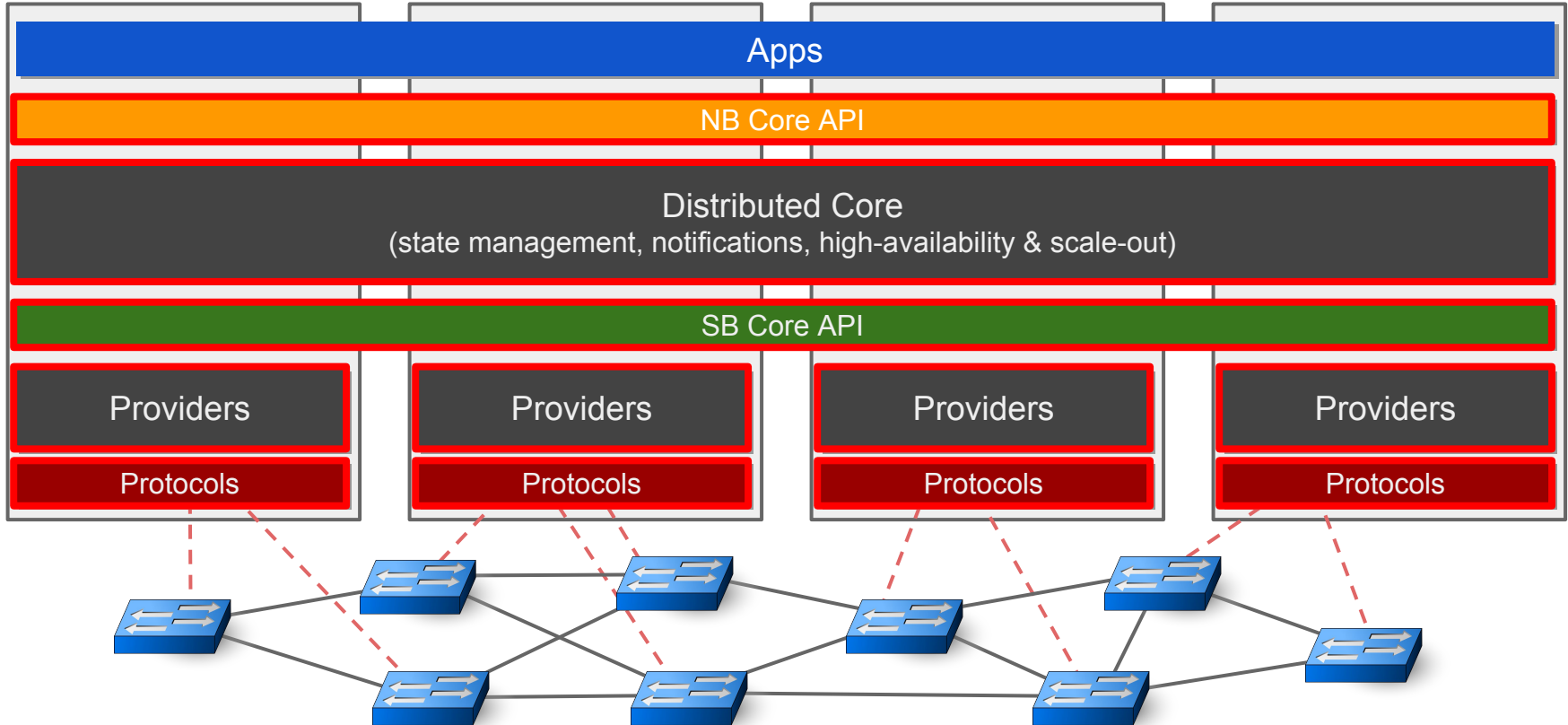


- High-availability, scalability and performance
 - required to sustain demands of service provider & enterprise networks
- Strong abstractions and simplicity
 - required for development of apps and solutions
- Protocol and device behaviour independence
 - avoid contouring and deformation due to protocol specifics
- Separation of concerns and modularity
 - allow tailoring and customization without speciating the code-base

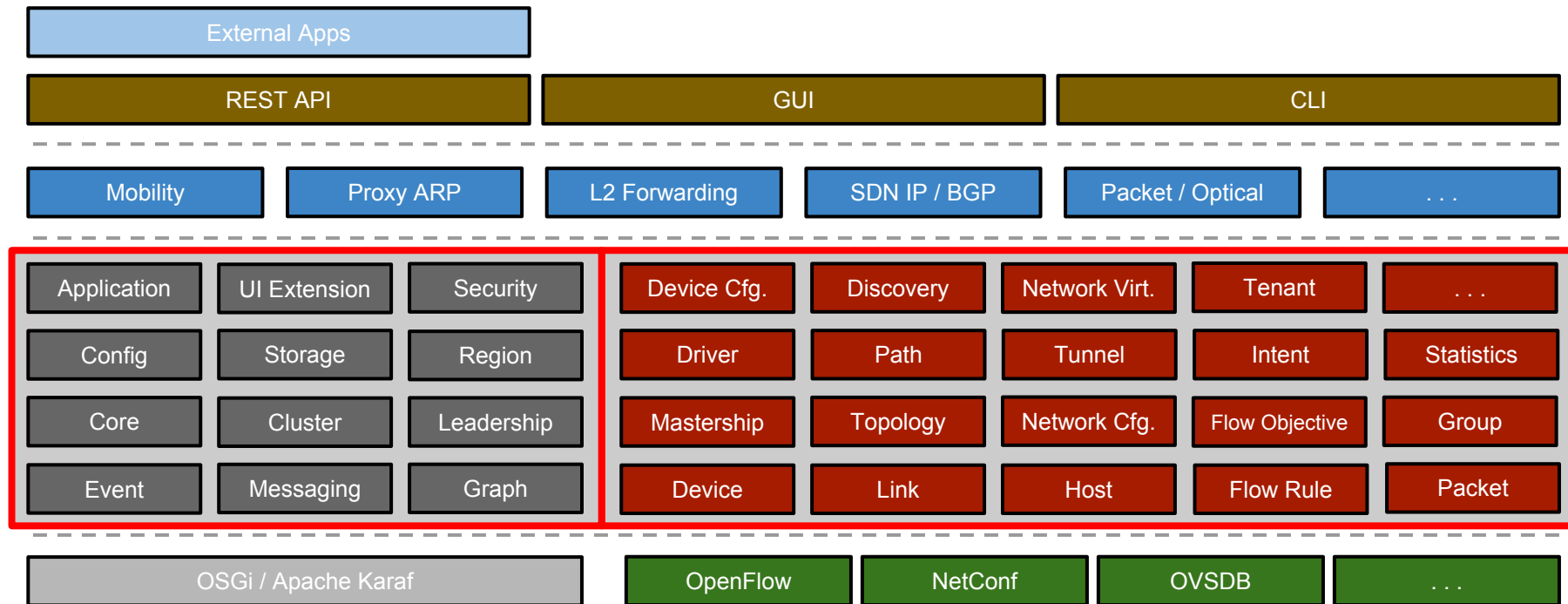
ONOS Distributed Architecture



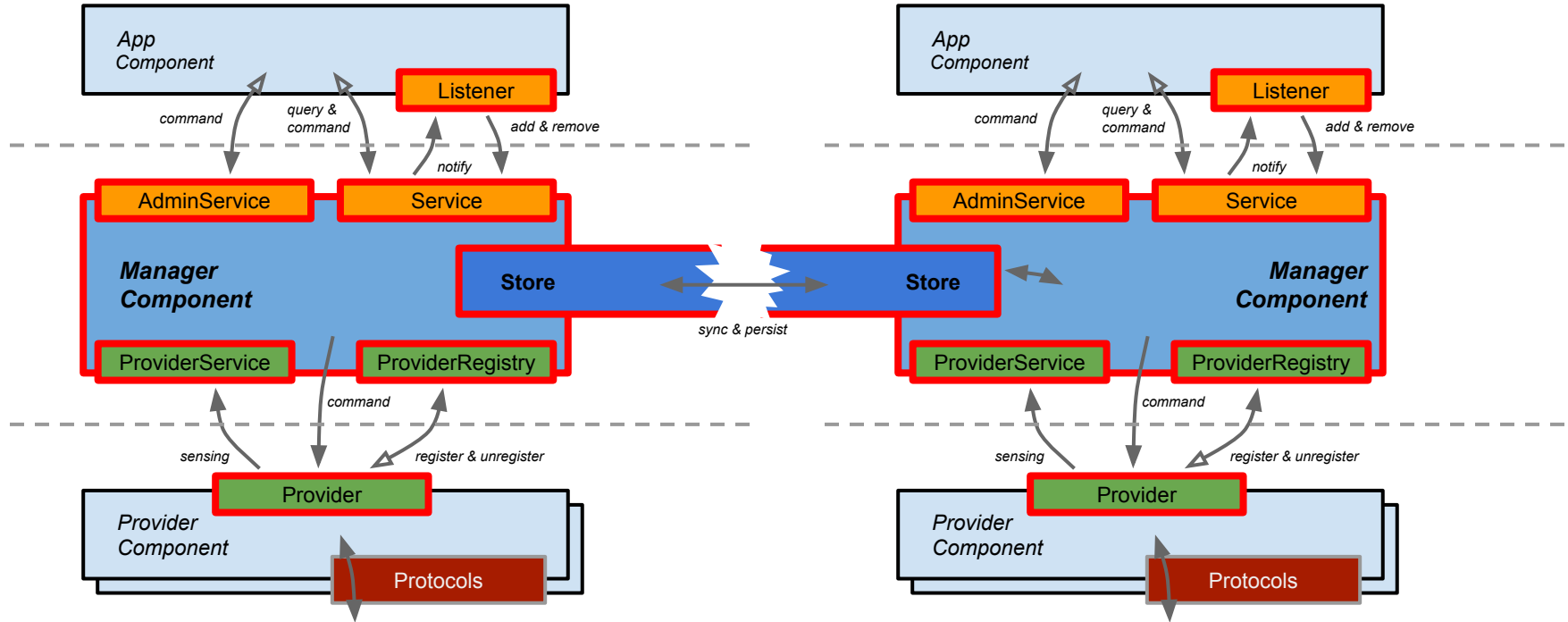
ONOS Distributed Architecture



ONOS Core Subsystems



ONOS Core Subsystem Structure



Key Northbound Abstractions

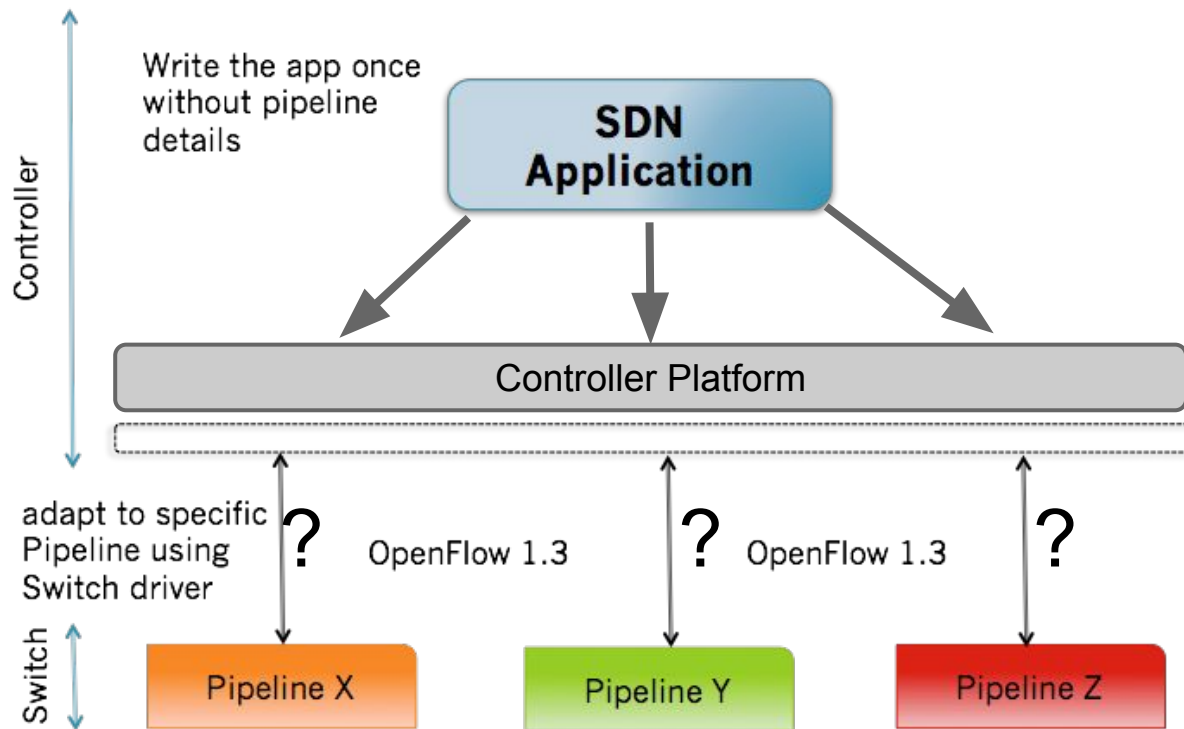


- Network Graph
 - Directed, cyclic graph comprising of infrastructure devices, infrastructure links and end-station hosts
- Flow Objective
 - Device-centric abstraction for programming data-plane flows in table pipeline-independent manner
- Intent
 - Network-centric abstraction for programming data-plane in topology-independent manner

Flow Objective Subsystem



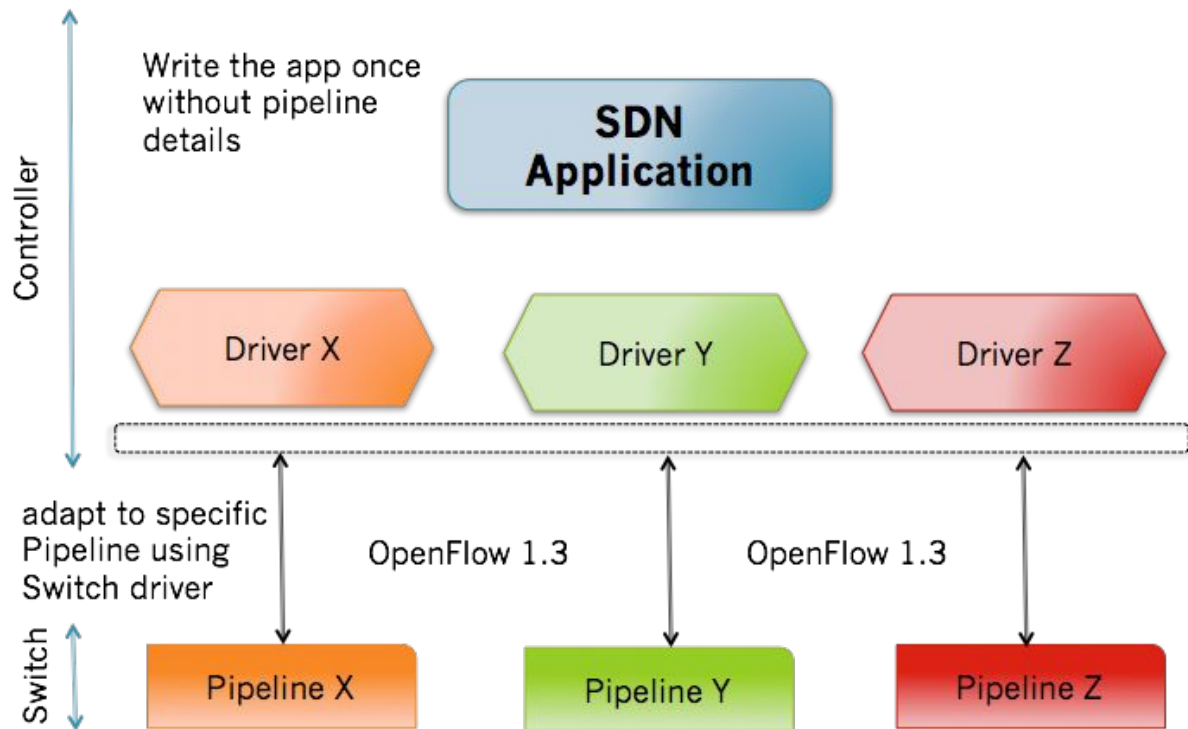
- **Problem:** Applications today must be pipeline aware, effectively making them applicable to specific HW.



Flow Objective Subsystem



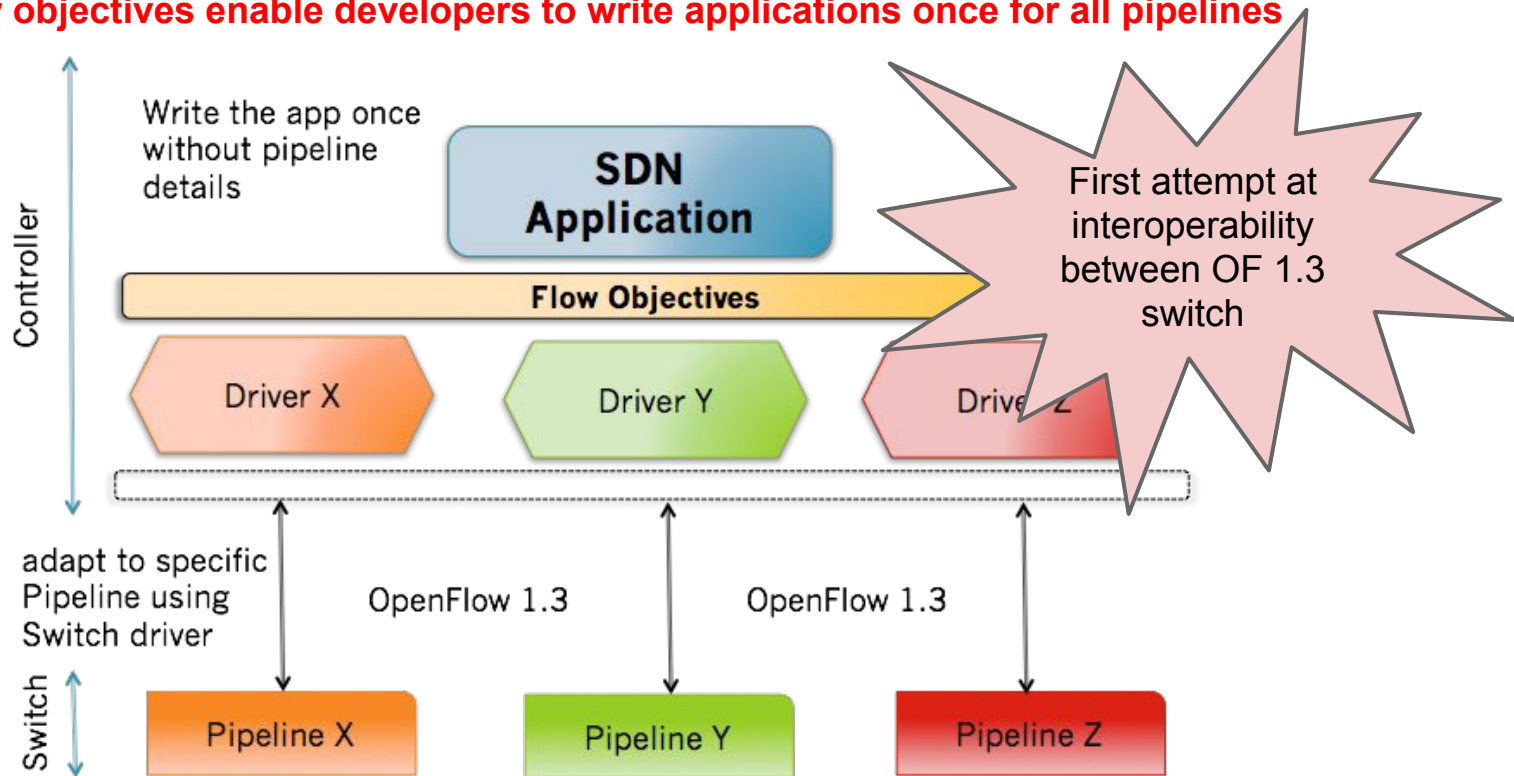
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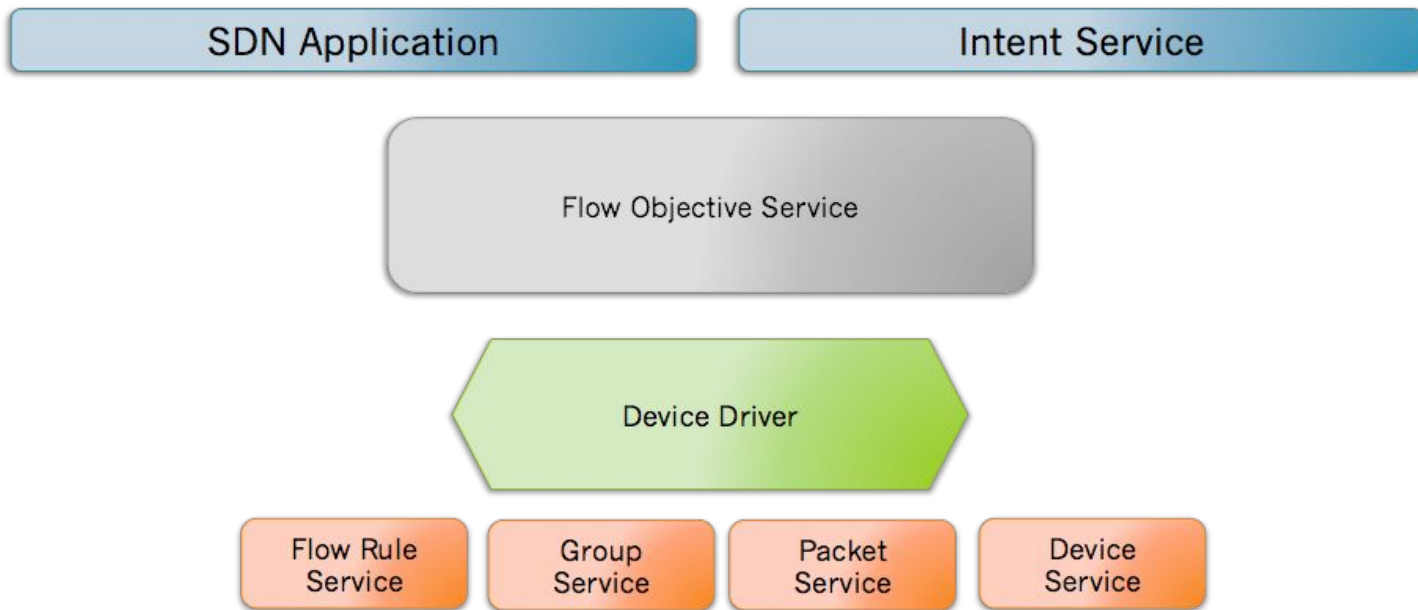
Flow Objective Abstraction



- **Problem:** Applications currently must be pipeline aware, effectively making applicable on specific HW.
Flow objectives enable developers to write applications once for all pipelines



Flow Objective Service



- Applications use Objective to take advantage multi-table architectures
- Other services also make use of the Objective service (eg. Intent Service)
- Device driver translates objectives to the specific flow rules for a given device

Flow Objectives



- Flow Objectives describe a SDN application's objective behind a **flow** it is sending to a **device**
- We currently only have three types of objectives:
 1. Filtering Objective
 2. Forwarding Objective
 3. Next Objective

Filtering Objective



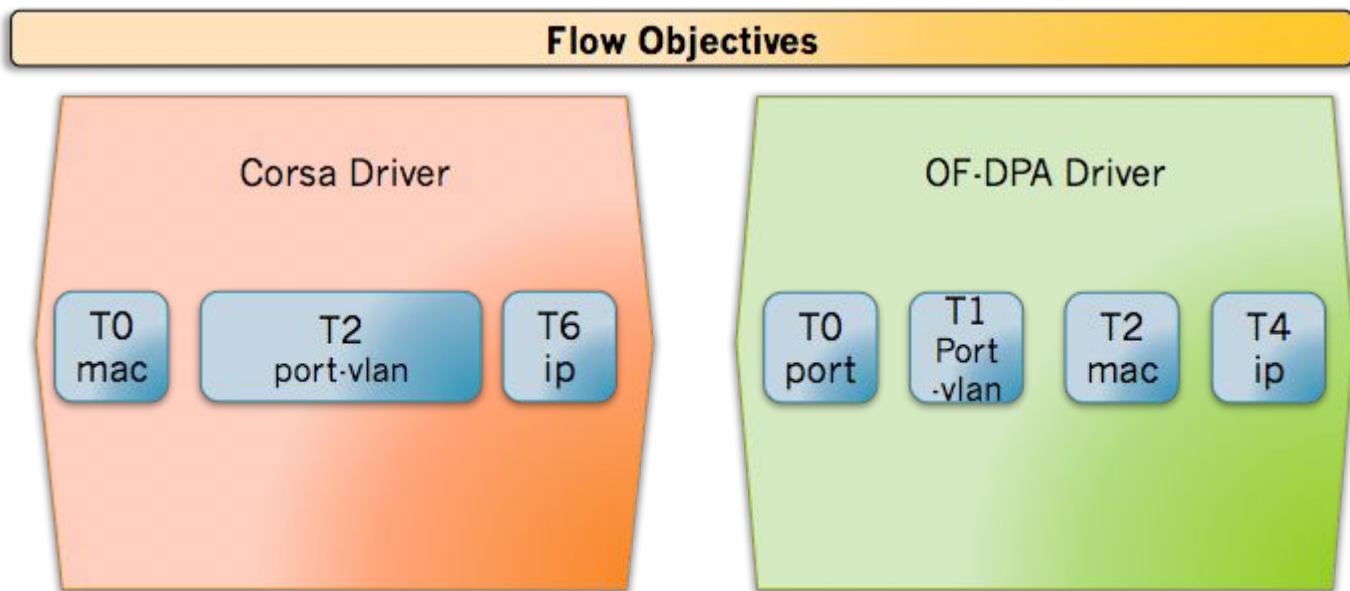
- ☐ Filter -> only Permit or Deny options
- ☐ On criteria (match fields)

Example:
Peering Router

Switch Port : **X**

Permit: MAC 1, VLAN 1, IP 1, 2, 3

Permit: MAC 1, VLAN 2, IP 4, 5



Filtering Objective



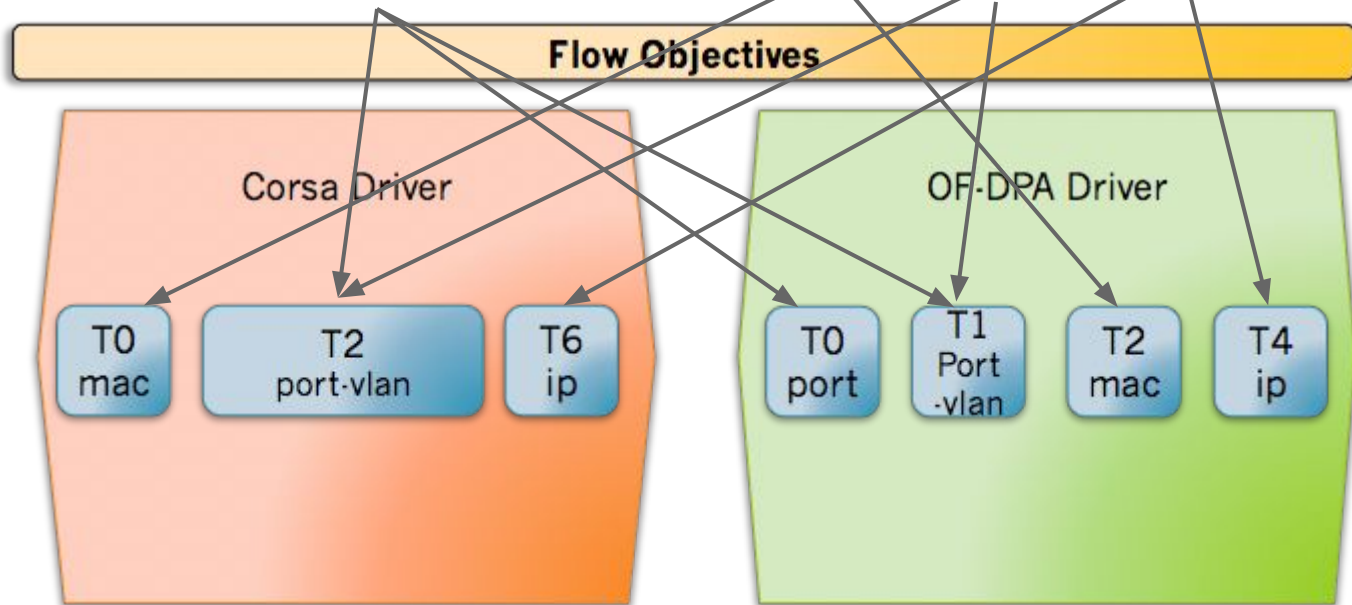
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Example:
Peering Router

Switch Port : X

Permit: MAC 1, VLAN 1, IP 1, 2, 3

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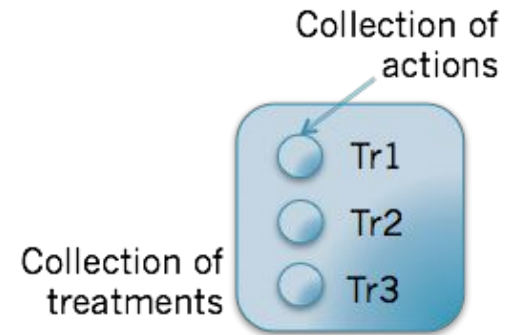
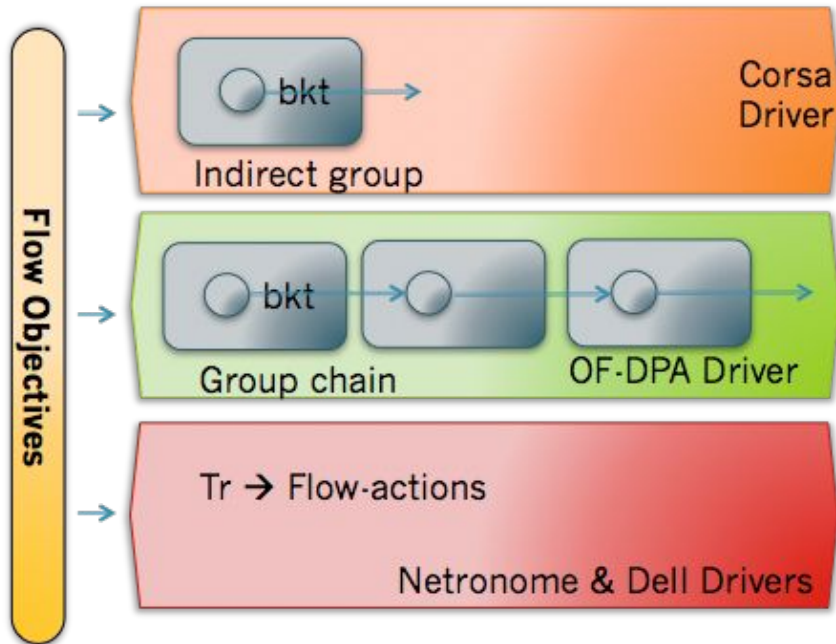


Next Objective

- Next -> next hop for forwarding
- Similar to OF group
- Keyed by a NextId used in Forwarding Objectives



Next Objective
Type: **SIMPLE**



Forwarding Objective



- Forwarding: { Selector -> Next Id }
- Forwarding Types: Specific or Versatile
 - Specific -> MAC, IP, MPLS forwarding tables
 - Versatile -> ACL table
- NextId is resolved to whatever the driver previously built for the corresponding Next Objective

Objectives - Simpler applications



```
for (InterfaceIpAddress ipAddr : intfIps) {  
    log.debug("adding rule for IPs: {}", ipAddr.ipAddress());  
    selector = DefaultTrafficSelector.builder();  
    treatment = DefaultTrafficTreatment.builder();
```

```
    selector.matchEthType(Ethernet.TYPE_IPV4);  
    selector.matchIPDst(IpPrefix.valueOf(ipAddr.ipAddress(), 32));  
    treatment.transition(Type.ACL);
```

```
    rule = new DefaultFlowRule(deviceId, selector.build(),  
                               treatment.build(), HIGHEST_PRIORITY, appId,  
                               0, true, FlowRule.Type.IP);
```

```
    ops = install ? ops.add(rule) : ops.remove(rule);  
}
```

```
for (MacAddress mac : intfMacs) {  
    log.debug("adding rule for MAC: {}", mac);  
    selector = DefaultTrafficSelector.builder();  
    treatment = DefaultTrafficTreatment.builder();
```

```
    selector.matchEthDst(mac);  
    treatment.transition(FlowRule.Type.VLAN_MPLS);
```

```
    rule = new DefaultFlowRule(deviceId, selector.build(),  
                               treatment.build(),  
                               CONTROLLER_PRIORITY, appId, 0,  
                               true, FlowRule.Type.FIRST);
```

```
    ops = install ? ops.add(rule) : ops.remove(rule);  
}
```

```
for (Map.Entry<PortNumber, VlanId> portVlan : portVlanPair.entrySet()) {  
    log.debug("adding rule for VLAN: {}", portVlan);  
    selector = DefaultTrafficSelector.builder();  
    treatment = DefaultTrafficTreatment.builder();
```

```
    selector.matchVlanId(portVlan.getValue());  
    selector.matchInPort(portVlan.getKey());  
    treatment.transition(Type.ETHER);  
    treatment.deferred().popVlan();
```

```
    rule = new DefaultFlowRule(deviceId, selector.build(),  
                               treatment.build(), CONTROLLER_PRIORITY, appId,  
                               0, true, FlowRule.Type.VLAN);
```

```
    ops = install ? ops.add(rule) : ops.remove(rule);  
}
```

```
private void processIntfFilters(boolean install, Set<Interface> intf) {  
    log.info("Processing {} router interfaces", intf.size());  
    for (Interface intf : intf) {  
        FilteringObjective.Builder fob = DefaultFilteringObjective.builder();  
        fob.withKey(Criteria.matchInPort(intf.connectPoint().port()))  
            .addCondition(Criteria.matchEthDst(intf.mac()))  
            .addCondition(Criteria.matchVlanId(intf.vlan()));  
        intf.ipAddresses().stream()  
            .forEach(ipaddr -> fob.addCondition(  
                Criteria.matchIPDst(ipaddr.subnetAddress())));  
        fob.permit().fromApp(appId);  
        flowObjectiveService.filter(deviceId,  
            Collections.singletonList(fob.add()));  
    }  
}
```

Flow Objective Summary



- *Flow Objective Service*: **Abstraction** for applications to be **pipeline unaware** while **benefiting** from scalable, multi-table architectures
- Aims to make it **simple** to write apps
- First attempt at achieving **interoperability** between OF 1.3 implementations

Building Network Applications



- Each application requires complex path computation and rule installation engines and state machines
- Inconsistent behavior in the face of failures
 - Failures may be handled in different ways (or not at all)
- Bugs need to be fixed in multiple places (applications)
- Expensive to upgrade/refactor behavior across all applications; e.g.
 - Improve performance
 - Support new types of devices
 - Implement better algorithms
- Difficult or impossible to resolve conflicts with other applications

Intent Framework



- Provides **high-level, network-centric** interface that focuses on *what* should be done rather than *how* it is specifically programmed
- Abstracts unnecessary network complexity from applications
- Maintains requested semantics as network changes
- High availability, scalability and high performance

Example Applications

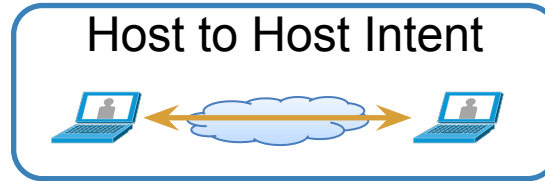


- **SDN-IP Peering**
 - Connect internal BGP software daemon to external BGP routers
 - Install learned routes to forward IP traffic to appropriate egress point
- **Multi-level (IP / Optical) Provisioning**
 - Provision optical paths/tunnels with constraints
- **Content Acquisition / Video Streaming (DirecTV)**
 - Establish multicast forwarding from a sender to set of receivers
- **Virtual Network Gateway (vBNG)**
 - Provide connectivity between a private host and the Internet
- **Bandwidth Calendaring**
 - Establish tunnels with bandwidth guarantees between two points at a given time

Intent Example



Intent Example



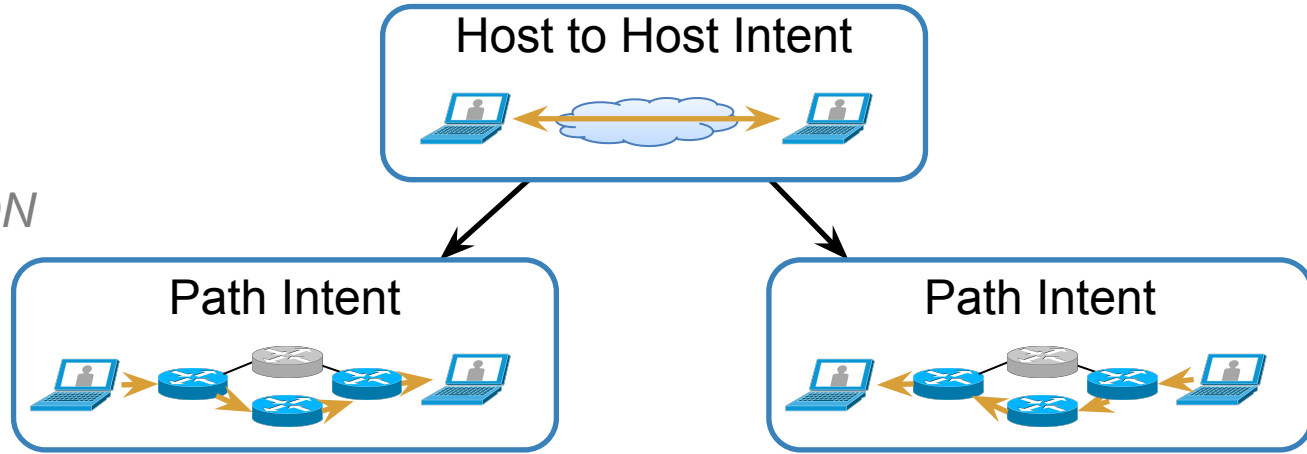
`submit()`

Intent Service API

Intent Example



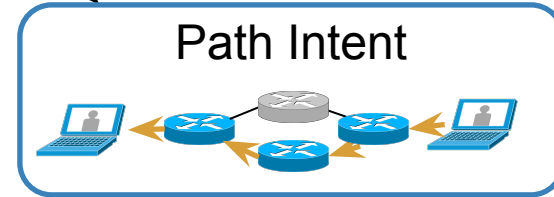
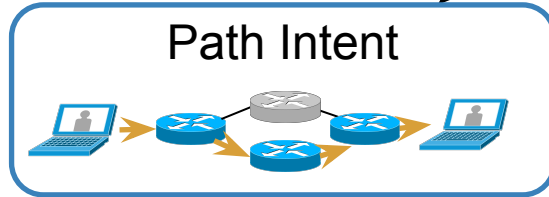
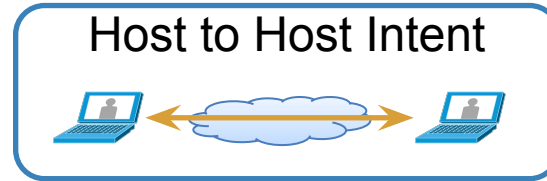
COMPILATION



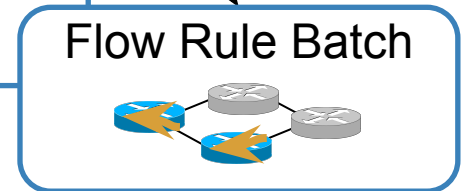
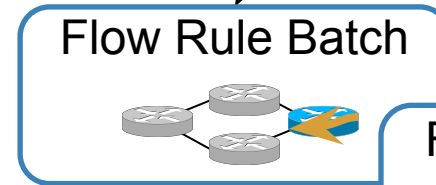
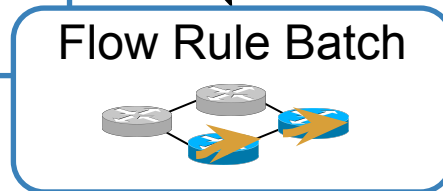
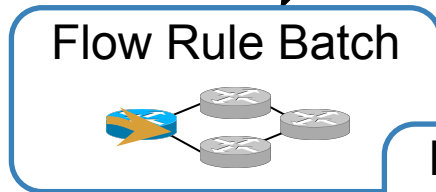
Intent Example



COMPILATION



INSTALLATION



Intent Framework Summary



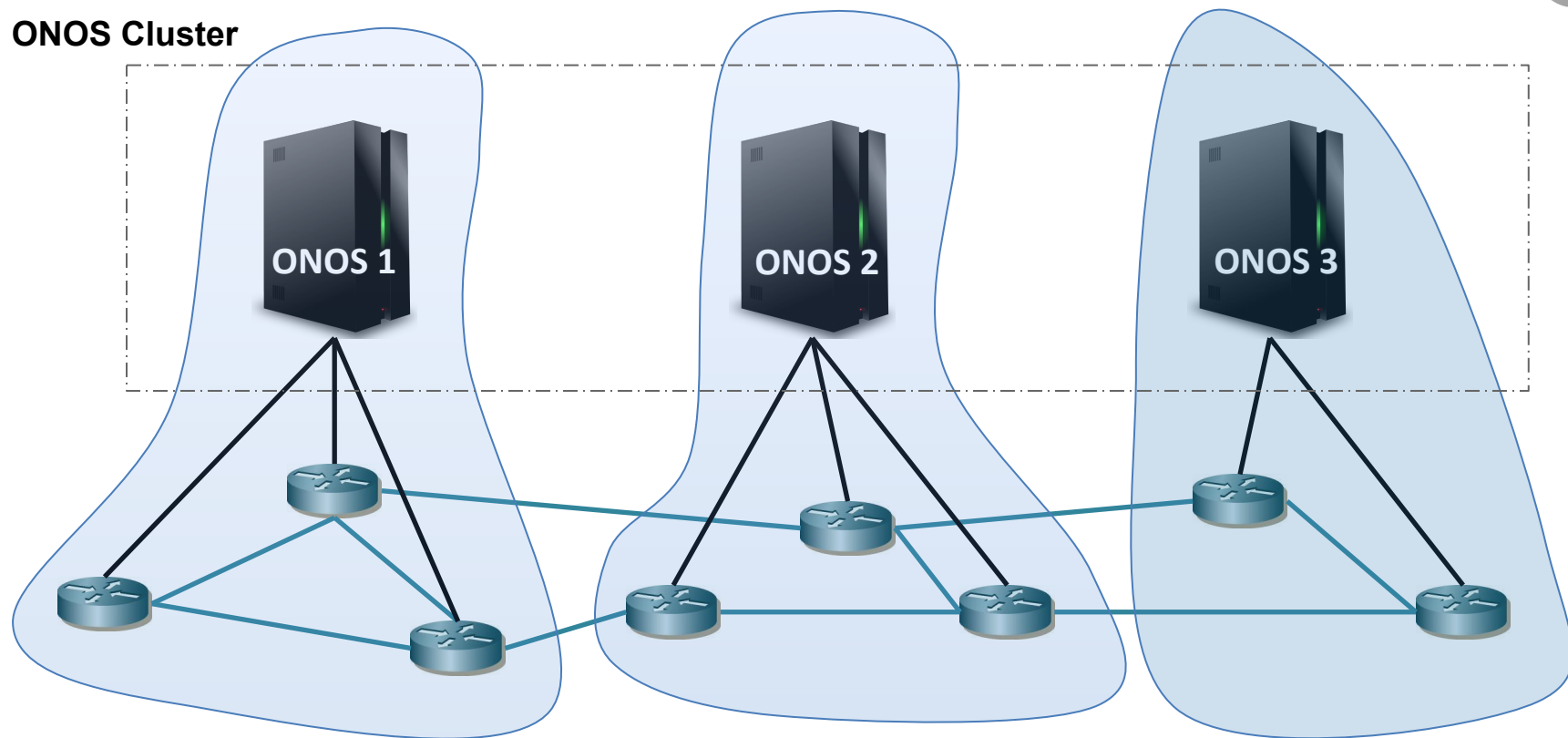
- Intents are a *network-centric programming abstraction* that **reduce application complexity**.
- Intents provide **device-agnostic behavior** with **persistence** and **high performance** across network failures.
- Intent framework has moved from prototype to **production** deployments.

ONOS Distributed Architecture



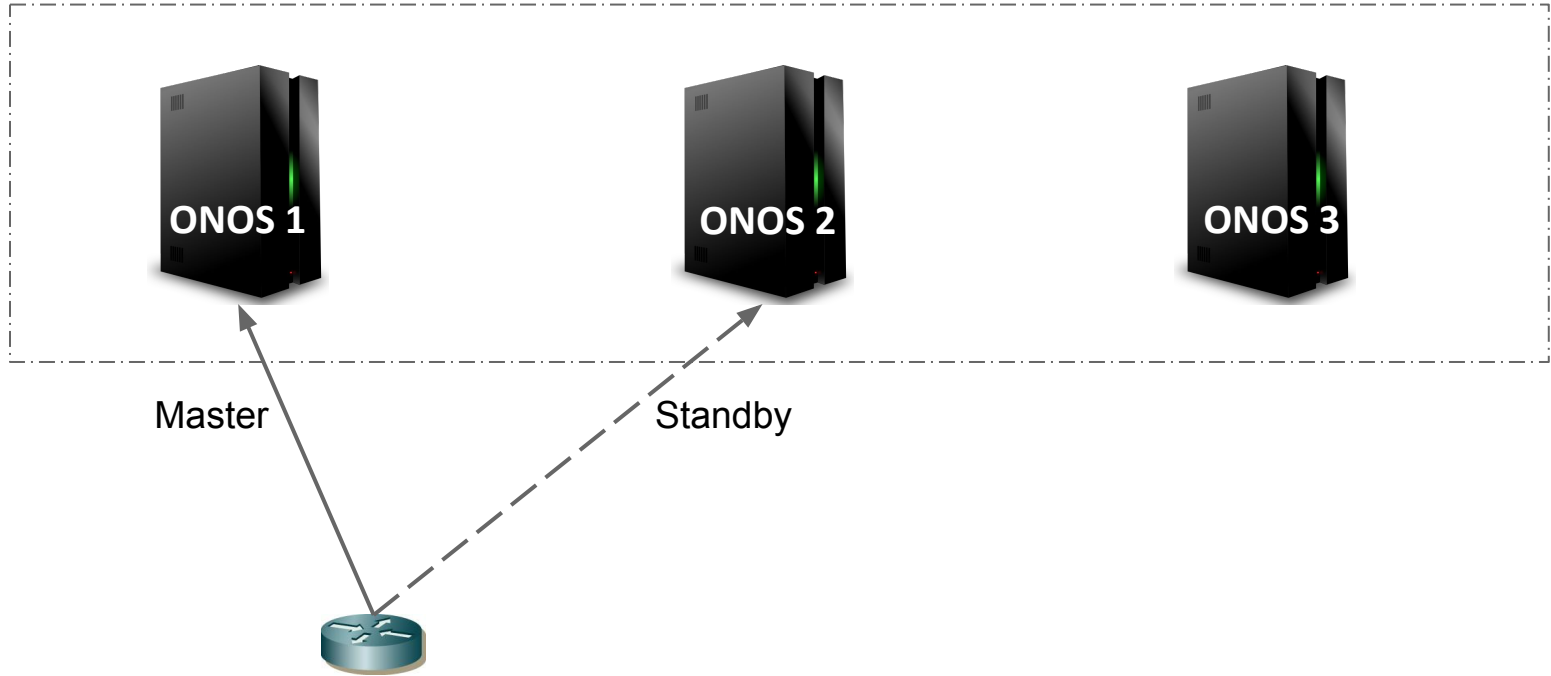
- Distributed
 - Set up as a cluster of instances
- Symmetric
 - Each instance runs identical software and configuration
- Fault-tolerant
 - Cluster remains operational in the face of node failures
- Location Transparent
 - A client can interact with any instance. The cluster presents the abstraction of a single logical instance
- Dynamic (*in progress*)
 - The cluster can be scaled up/down to meet usage demands

ONOS Cluster

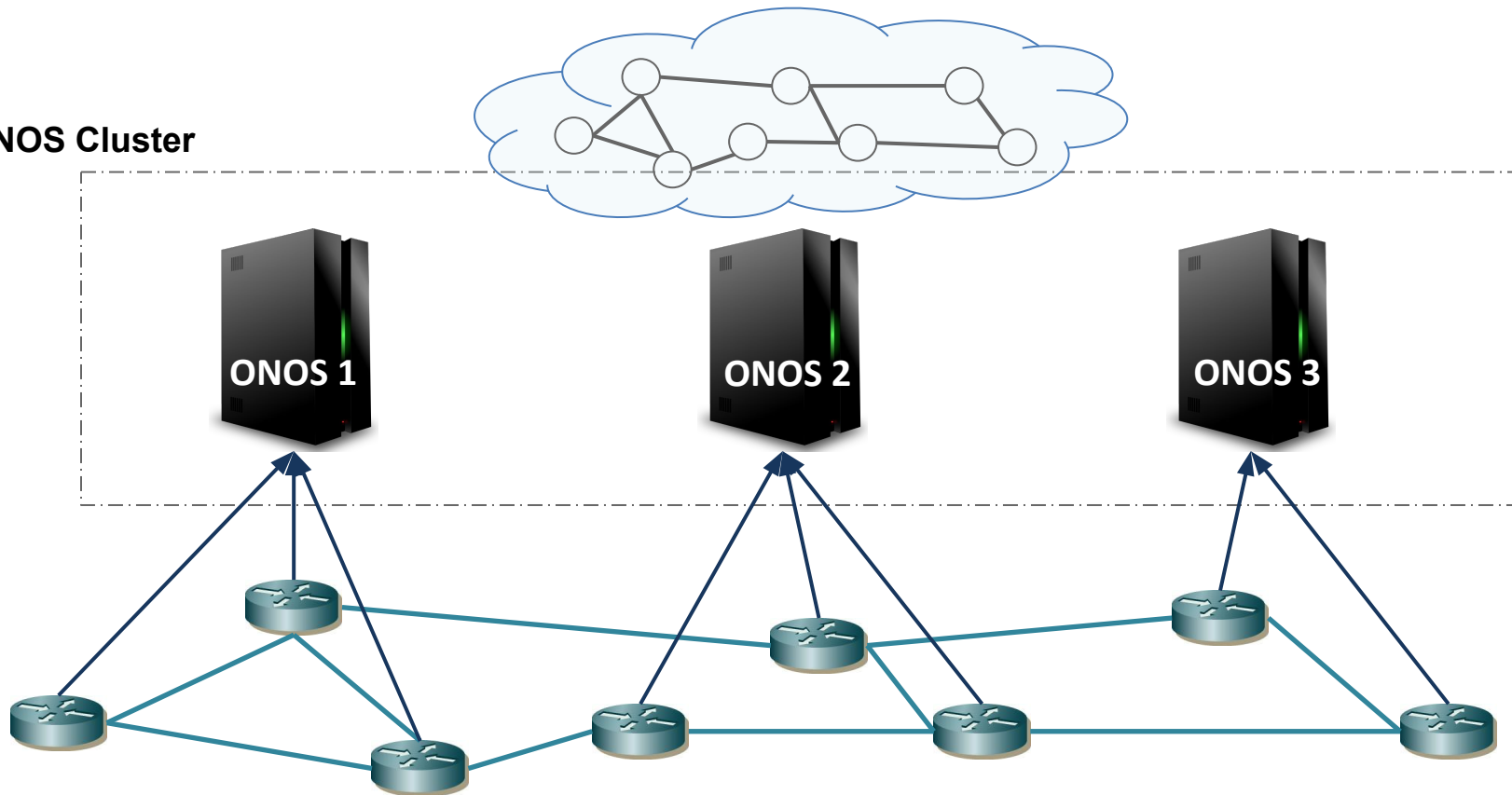


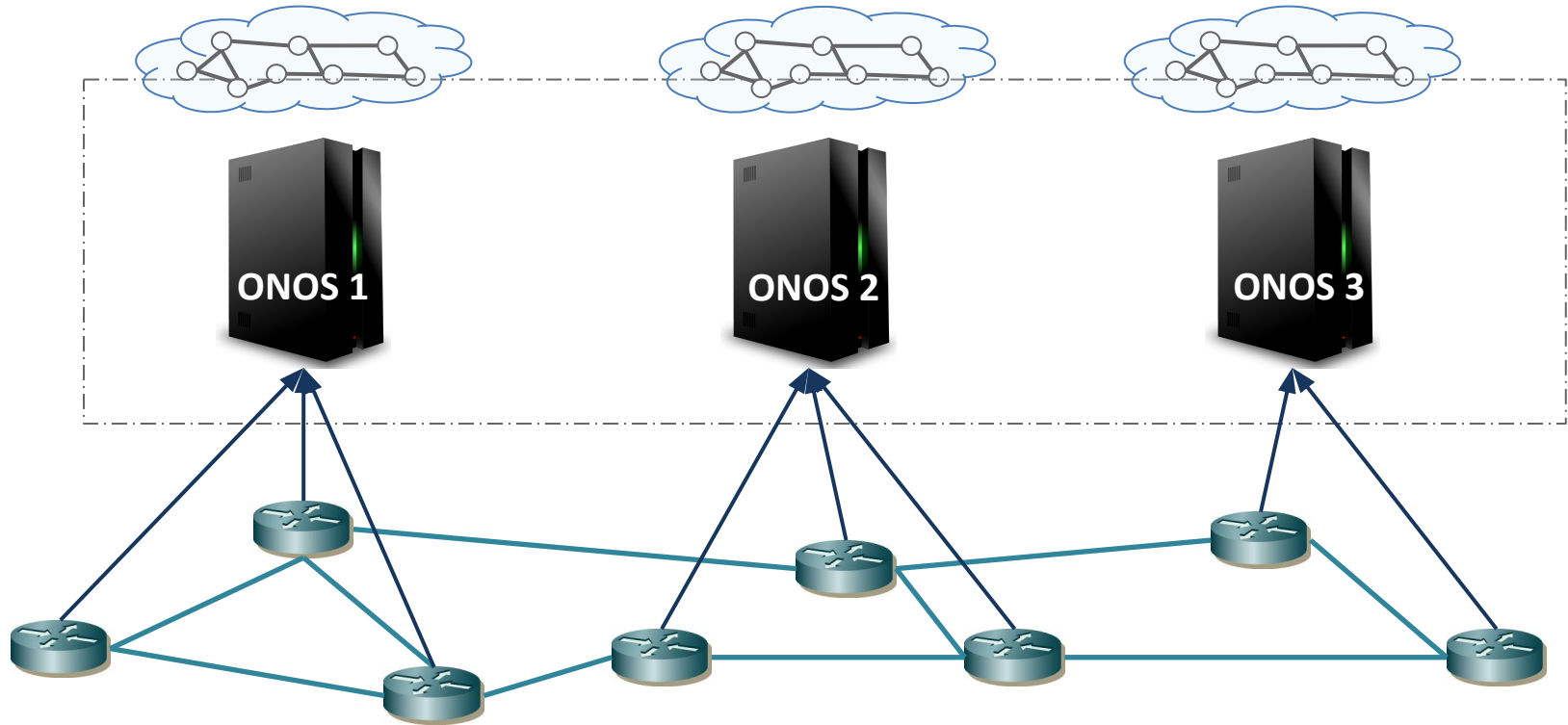


ONOS Cluster

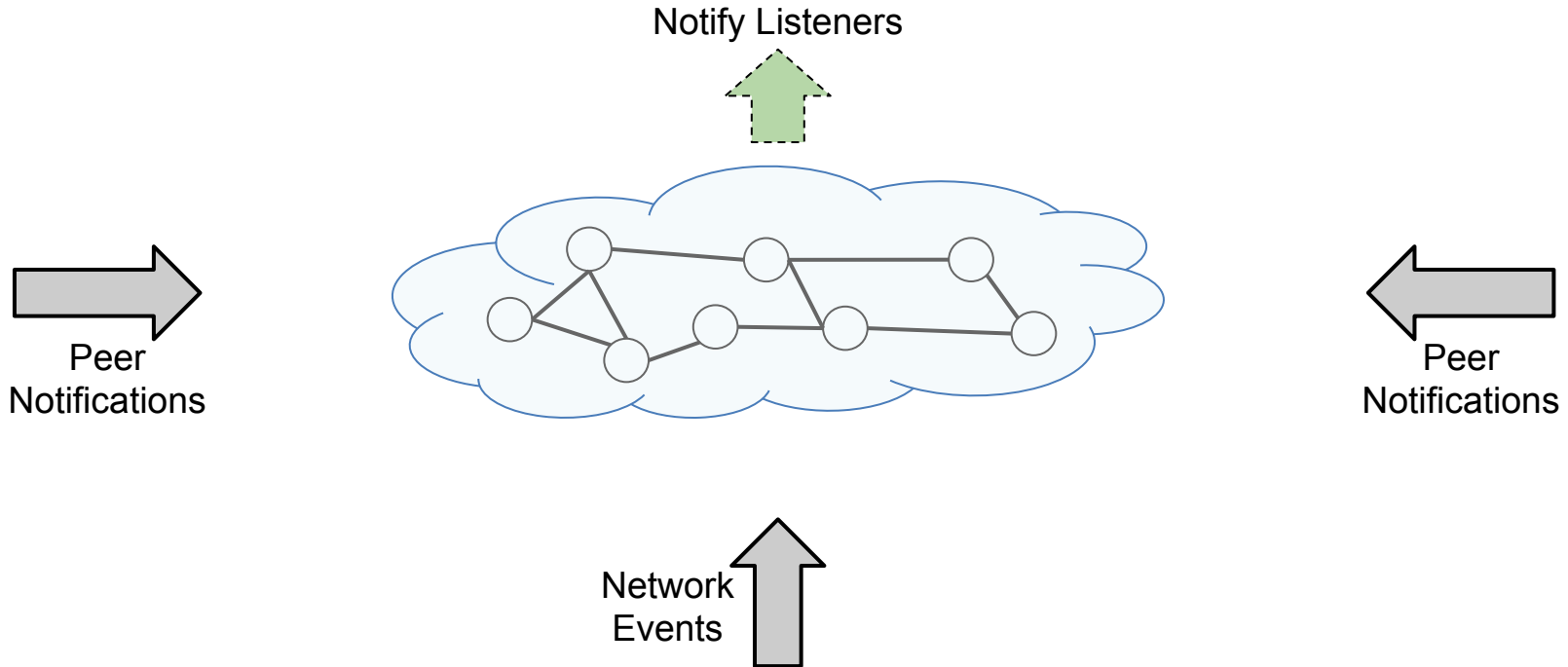


ONOS Cluster





Topology state inside a Node



Network Events and Ordering



Network Events are timestamped with (t, s)

$t \rightarrow$ mastership term number

$s \rightarrow$ sequence number in term

Series of timestamps for port X:

... (4, 4) (4, 5) (5, 1) (5, 2) ...



mastership term boundary

Network Topology State



- Eventually Consistent: Reads are **monotonically consistent**
- Low overhead reads and writes
 - 2-3 ms latency for reacting to network events
- Gossip based Anti-Entropy protocol fixes divergent copies
- Generalized as **EventuallyConsistentMap<K, V>**

State Management in ONOS



- Core platform feature
- Typically one of hardest pieces to get right and it is better to solve it once
- Better if applications can focus on business logic
- ONOS exposes a set of primitives to cater to different use cases
- Primitives span the consistency continuum

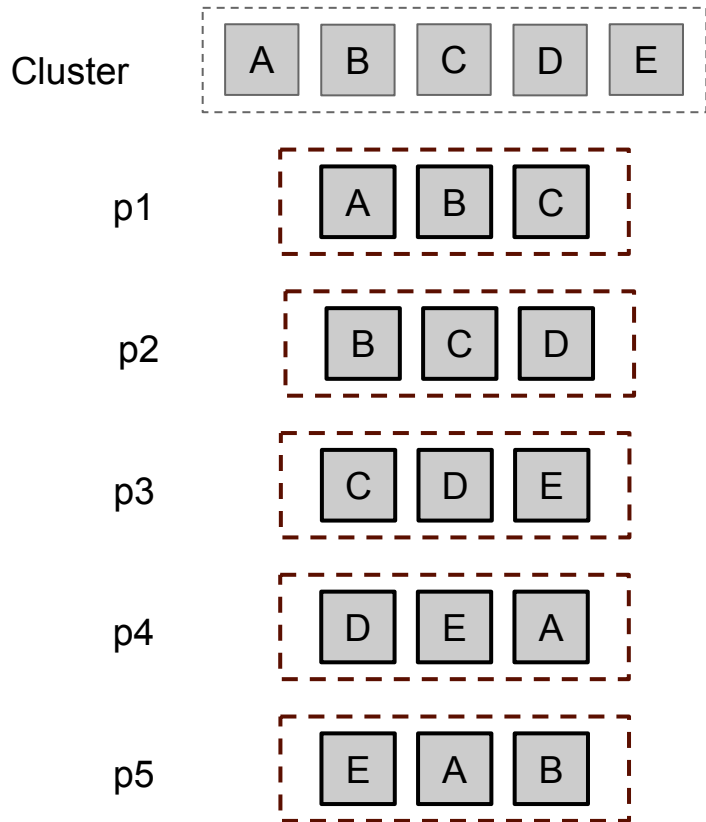


ONOS Distributed Primitives



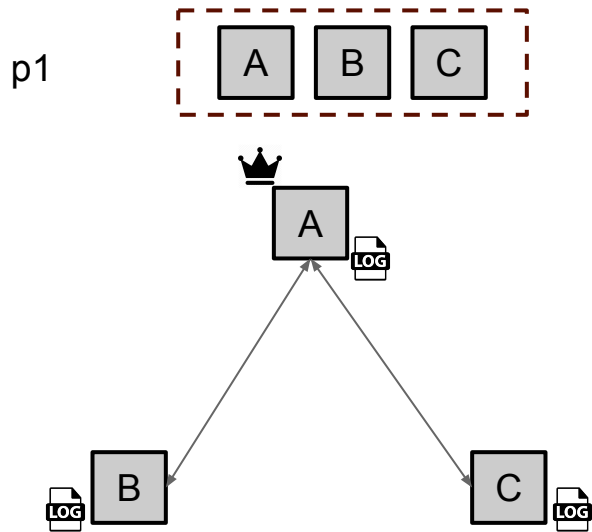
- **EventuallyConsistentMap<K, V>**
 - Map abstraction with eventual consistency guarantee
- **ConsistentMap<K, V>**
 - Map abstraction with strong linearizable consistency
- **LeadershipService**
 - Distributed Locking primitive
- **DistributedQueue<E>**
 - Distributed FIFO queue with long poll support
- **DistributedSet<E>**
 - Distributed collection of unique elements
- **AtomicCounter**
 - Distributed version of Java AtomicLong
- **AtomicValue<V>**
 - Distributed version of Java AtomicReference

Behind the scenes...



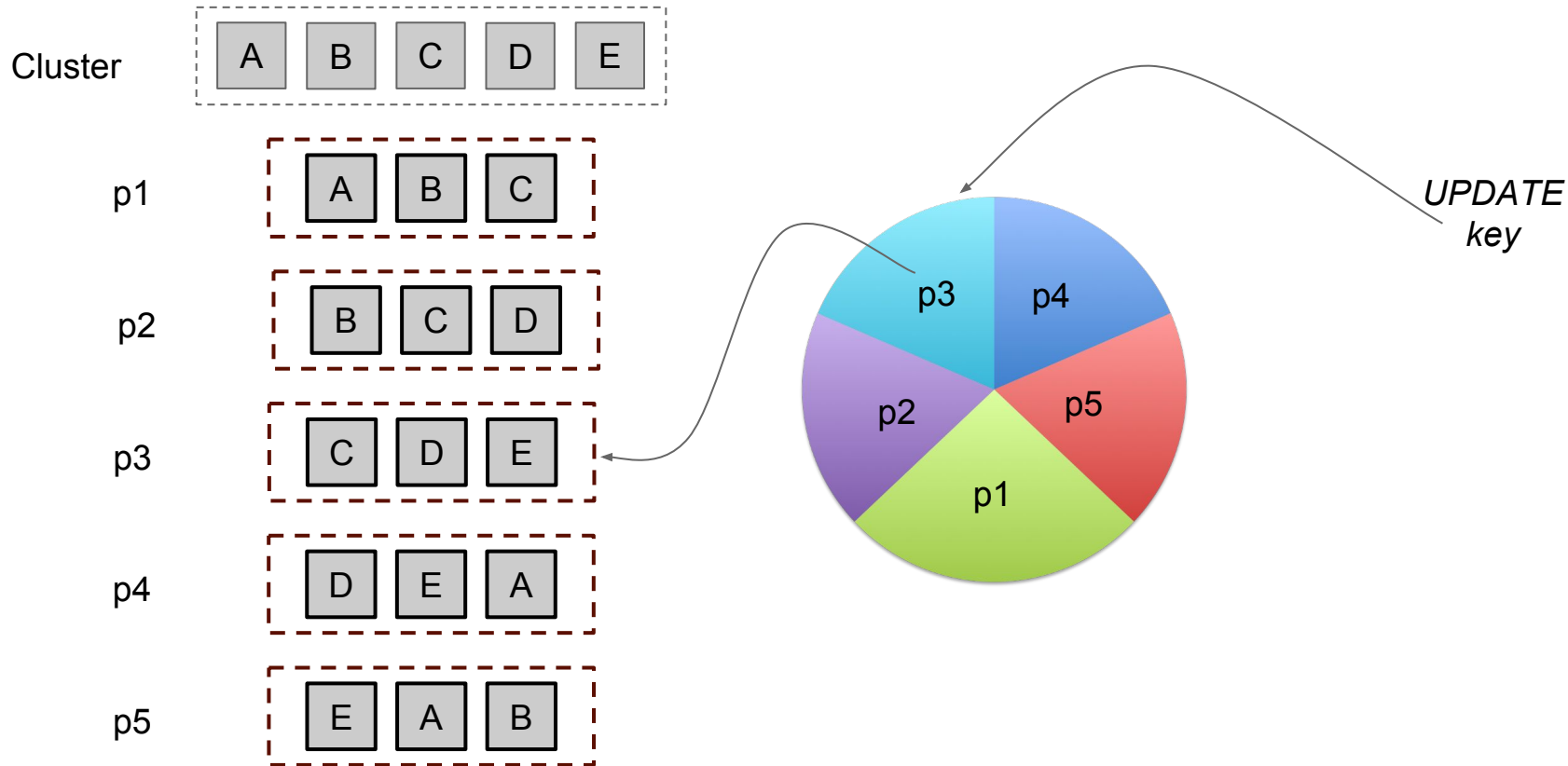
Data is partitioned into
Replica Sets

Inside a Replica Set

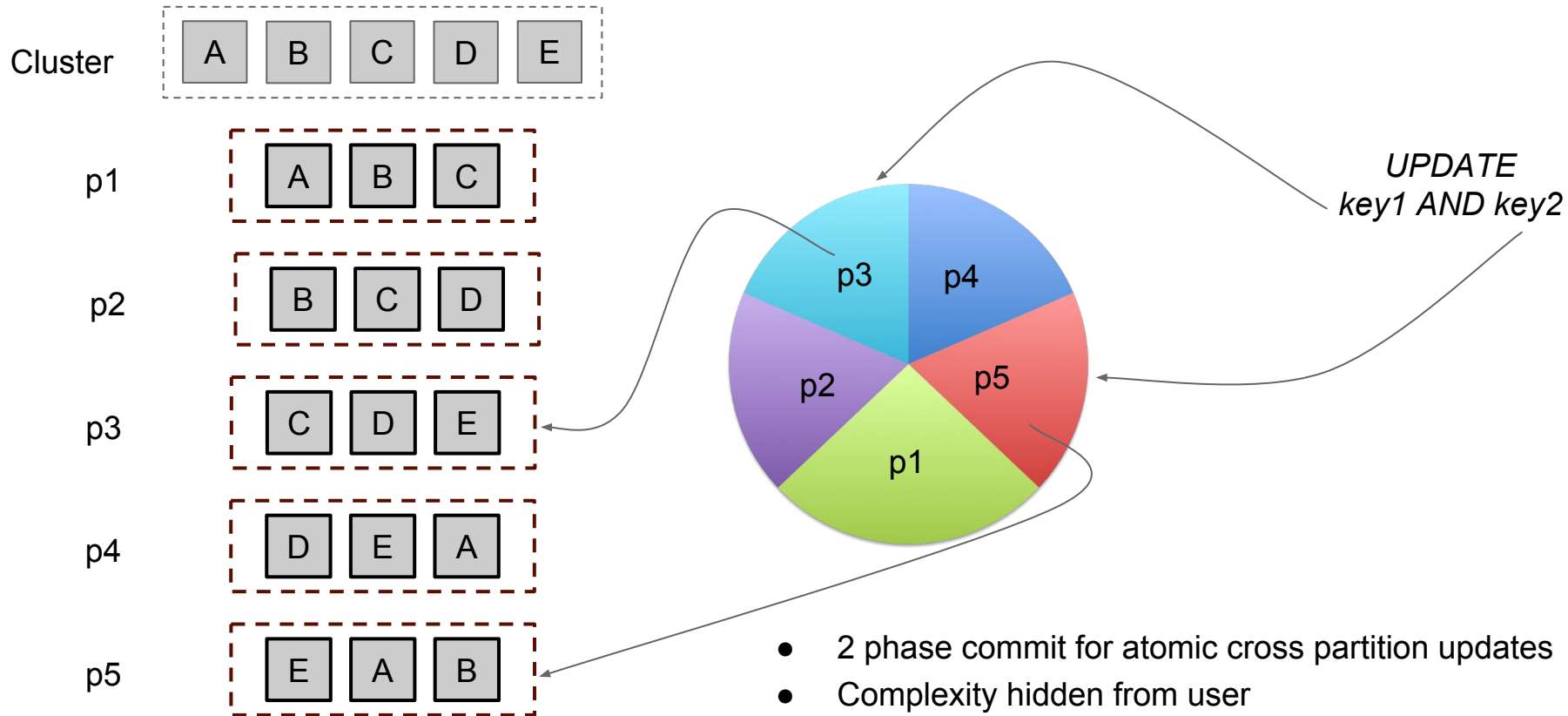


Raft consensus is used to maintain a Replicated State Machine

Data placement



Transactional Updates



Dynamic Clustering



- Ability to grow/shrink a cluster to suit usage demands
- Extract cluster metadata to a separate logical store
- Reshuffle data and control responsibilities to ensure fault-tolerance and load balance

Performance Metrics



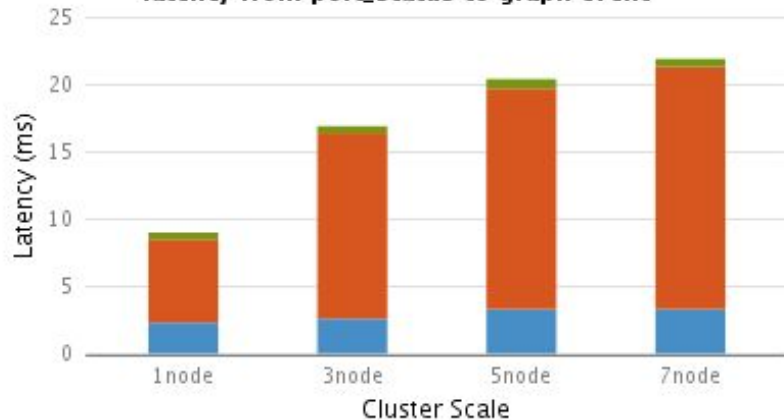
- Device & link sensing latency
 - measure how fast can controller react to environment changes, such as switch or port down to rebuild the network graph and notify apps
- Flow rule operations throughput
 - measure how many flow rule operations can be issued against the controller and characterize relationship of throughput with cluster size
- Intent operations throughput
 - measure how many intent operations can be issued against controller cluster and characterize relationship of throughput with cluster size
- Intent operations latency
 - measure how fast can the controller react to environment changes and reprovision intents on the data-plane and characterize scalability

Link Up/Down Latency



Link Up Latency Tests (Mean)

latency from port_status to graph event



■ OFP Port Status -> Device Event ■ Device Event -> Link Event ■ Link Event -> Graph Event

- Since we use LLDP & BDDP to discover links, it takes longer to discover a link coming up than going down

Link Down Latency Tests (Mean)

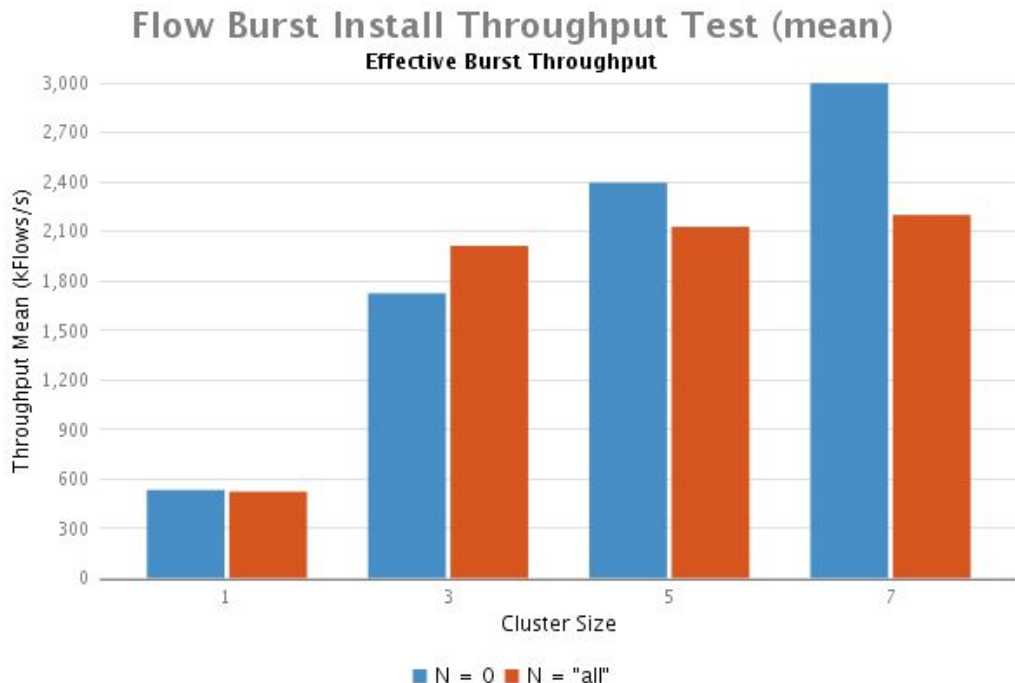
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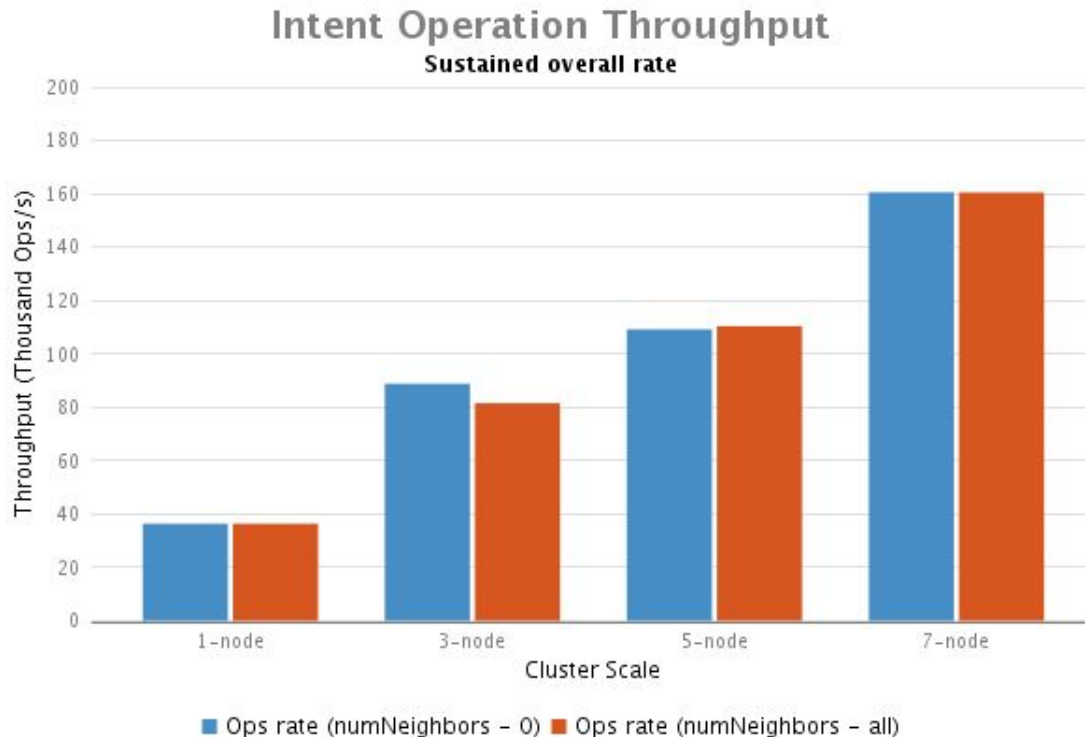
- Port down event trigger immediate teardown of the link.

Flow Throughput results



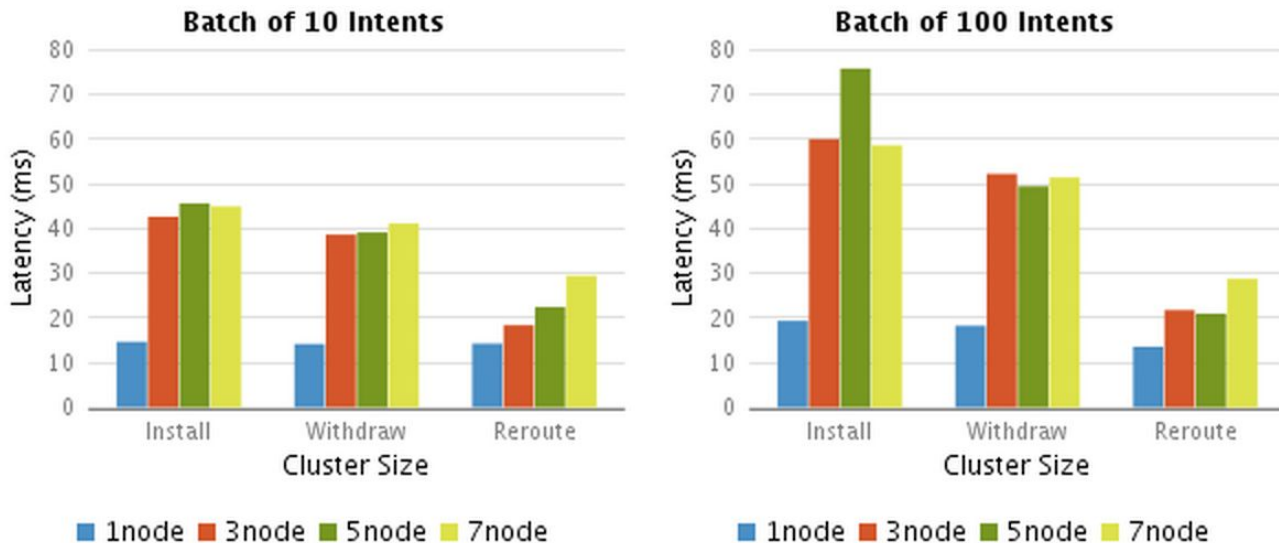
- Single instance can install over 500K flows per second
- ONOS can handle 3M local and 2M non local flow installations
- With 1-3 ONOS instances, the flow setup rate remains constant no matter how many neighbours are involved
- With more than 3 instances injecting load the flow performance drops off due to extra coordination required.

Intent Throughput Results



- Processing clearly scales as cluster size increases

Intent Latency Results



- Less than 100ms to install or withdraw a batch of intents
- Less than 50ms to process and react to network events
 - Slightly faster because intent objects are already replicated



Software Defined Transformation of Service Provider Networks

Join the journey @ onosproject.org